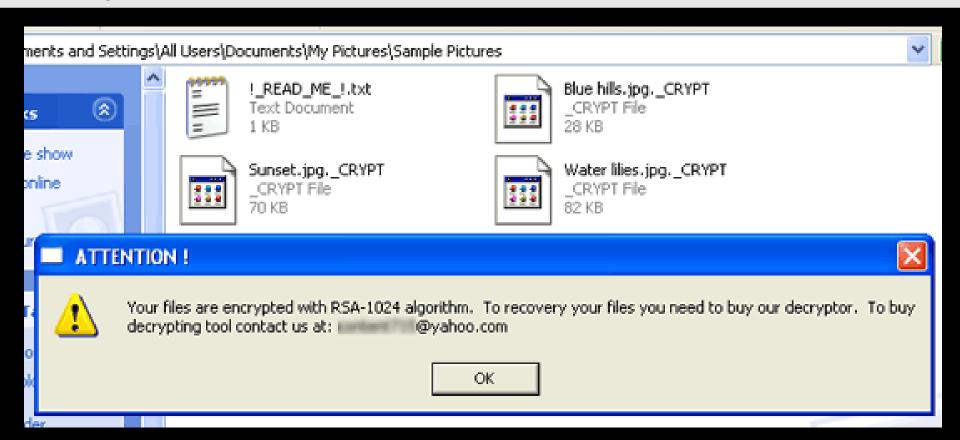
Cryptanalysis of the Gpcode.ak ransomware virus

Eran Tromer

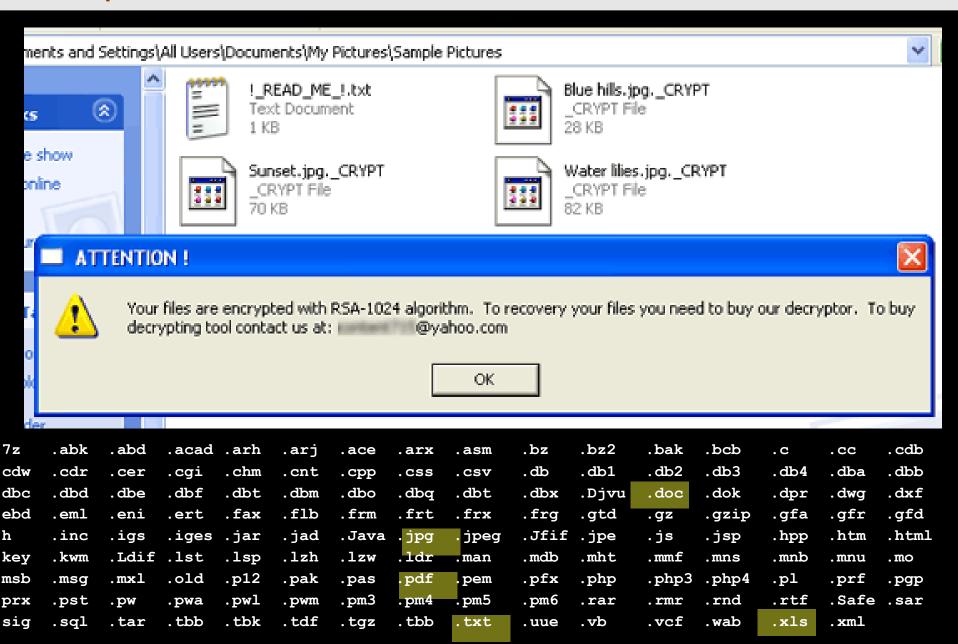
MIT



The Gpcode.ak ransomware virus



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Email

Next, you should send \$100 to Liberty Reserve account U6890784 or E-Gold account 5431725 (www.e-gold.com) To buy E-currency you may use exchange service, see or any other. In the transfer description specify your e-mail. After receive your payment, we send decryptor to your e-mail. For check our guarantee you may send us one any encrypted file (with cipher key, specified in any !_READ_ME_!.txt file, being in the directorys with the encrypted files). We decrypt it and send to you originally decrypted file.

Best Regards, Daniel Robertson



Email 2

The price of decryptor is 200 USD. For payment you may use one of following variants: 1. Payment to E-Gold account 5437838 (www.e-gold.com). 2. Payment to Liberty Reserve account U6890784 (www.libertyreserve.com). 3. If you do not make one of this variants, contact us for decision it. For check our guarantee you may send us ONE any encrypted file. We decrypt it and send to you originally decrypted file. For any questions contact us via e-mail.

Best regards. Paul Dyke

(Russian criminals via a Chinese ISP)



Reverse-engineering Gpcode.ak

(based on a sample from Kaspersky Lab)

Virus payload

- Generate a random RC4 machine key $K_{
 m M}$
- For every file *f* :
 - -Generate random file nonce N_f
 - -Derive an RC4 file key K_f from N_i and $K_{
 m M}$
 - -Encrypt the file using K_f and prepend N_f to the ciphertext
 - Delete original file
- Encrypt $K_{\rm M}$ under an embedded 1024-bit RSA public key and write it to READ ME !.txt
- Forget $K_{
 m M}$

All in 8030 bytes. Uses Windows CryptoAPI.



Decryption

 Victim sends _READ_ME_!.txt and \$200 to "Daniel Robertson".

• "Daniel" decrypts the machine key $K_{
m M}$ using

his RSA private key and sends it to the victim, embedded in a "decryptor".

 Decryptor rederives the file keys.





Cryptanalysis

The Gpcode.ak 1024-bit RSA Factoring Challenge

- *e*=65537
- $\begin{array}{l} \bullet \quad n_1 = \text{c}0\text{c}21\text{d}693223\text{d}68\text{f}b573\text{c}5318982595799\text{d}2d295\text{e}d37\text{d}a38\text{b}e41\text{a}c848\\ 6\text{e}f900\text{a}ee78\text{b}4729668\text{f}c920\text{e}e15\text{f}e0\text{b}587\text{d}1\text{b}61894\text{d}1\text{e}e15\text{f}5793\text{c}18\text{e}2\text{d}2\\ \text{c}8\text{c}c64\text{b}0539\text{e}01\text{d}088\text{e}41\text{e}0\text{e}\text{a}\text{f}d85055\text{b}6\text{f}55\text{d}232749\text{e}\text{f}48\text{c}\text{f}\text{e}6\text{f}\text{e}905011\text{c}1\\ 97\text{e}4\text{a}c6498\text{c}0\text{e}60567819\text{e}\text{a}\text{b}1471\text{c}\text{f}\text{a}4\text{f}2\text{f}4\text{a}27\text{e}3275\text{b}62\text{d}4\text{d}1\text{b}\text{f}0\text{c}79\text{c}665\\ 46782\text{b}81\text{e}93\text{f}85\text{d}_{16} \end{array}$

or

 $\begin{array}{l} n_2 = & \text{d}6046 \text{a} \\ \text{d}612773 \text{d}f8 \\ \text{d}c98 \\ \text{b}4033 \\ \text{a}3205 \\ \text{f}21c44703 \\ \text{d}a73 \\ \text{d}91631 \\ \text{c}6523 \\ \text{f}e735 \\ \text{d}607247 \\ \text{c}c9 \\ \text{a}5e0 \\ \text{f}936 \\ \text{e}d75 \\ \text{c}75 \\ \text{a}c7 \\ \text{c}e5 \\ \text{c}6e \\ \text{f}32 \\ \text{f}ff996 \\ \text{e}94 \\ \text{c}01e \\ \text{d}301289479 \\ \text{d}8d708 \\ \text{b}2c030 \\ \text{f}b79 \\ \text{d}225 \\ \text{a}7e0 \\ \text{b}e2 \\ \text{a}64e5 \\ \text{e}46e8336 \\ \text{e}03e0 \\ \text{f}6ced482939 \\ \text{f}c5715 \\ \text{1}4b8 \\ \text{d}7280 \\ \text{a}b5 \\ \text{f}4045106 \\ \text{b}7a4b7 \\ \text{f}a6b \\ \text{d}586c8 \\ \text{d}26d \\ \text{a}fb14b3 \\ \text{d}e71ca521432 \\ \text{d}65d \\ \text{d}538526 \\ \text{f}308 \\ \text{a}fb_{16} \\ \text{e}3086 \\ \text{f}3086 \\ \text{f}3086$



Cost of factoring a 1024-bit RSA modulus

- Exhaustive search
 - Don't be silly
- Number Field Sieve
 - TWIRL (\$1M x year)
 - SHARK (\$200M x year)
 - + matrix step
 - + NRE
 - + energy



Key derivation weakness (there are others)

$$K_f \leftarrow \mathrm{RC4}_{K_{\mathrm{M}}}(N_f)$$

But RC4 is a stream cipher!

$$K_f = K'_{\mathrm{M}} \oplus N_f$$
 where $K'_{\mathrm{M}} = \overline{\mathrm{RC4}}(K_{\mathrm{M}})$

Every file with known plaintext header gives:

$$N_f, \ \overline{\mathrm{RC4}}(K'_{\mathrm{M}} \oplus N_f)$$

Goal: recover the effective machine key K'_{M} .



Effective machine key recovery attack

- The setting is analogous to the [Fluhrer Mantin Shamir 2001] WEP attack, but the IV is XORed instead of prepended.
- Exploits (different) statistical biases in RC4:
 - [Roos 1995] [Maitra Paul 2007]
 - Tews Weinmann Pyshkin 2007]
 - [Klein 2008]– [Paul Rathi Maitra 2008]
- Combine biases optimally using Bayesian inference
- Key recovery using ~10000 encrypted files (vs. ~25000 packets for WEP)



Conclusions

- Cryptanalysis saves the day!
 (if you have have a few thousand .jpg files on your disk)
- Open problem: UC-secure ransom scheme in the standard model

